

Variations on Pebbling and Graham's Conjecture

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Talk structure

Pebbling numbers

Products and Graham's Conjecture

Variations

- Optimal pebbling

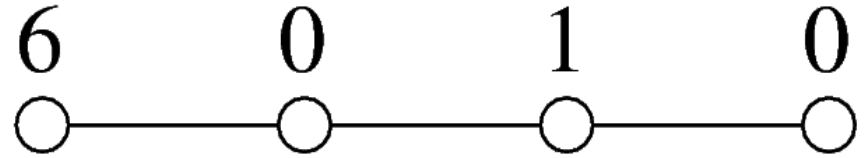
- Weighted graphs

- Choosing target distributions

Basic notions

Distributions on G :

$$D: V(G) \rightarrow \mathbb{N}$$



$D(v)$ counts pebbles
on v

Pebbling moves

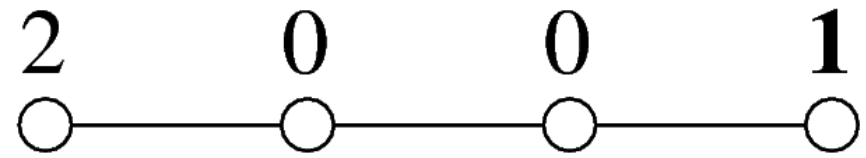
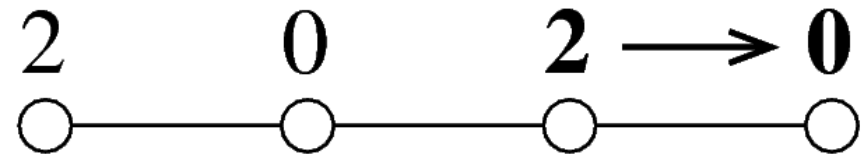
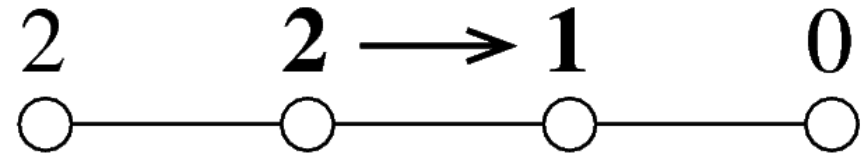
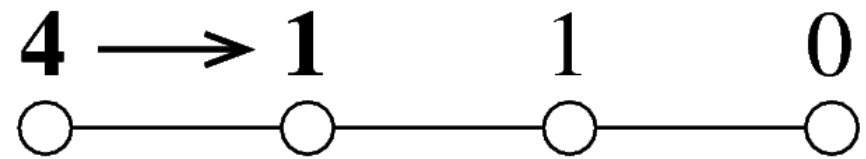
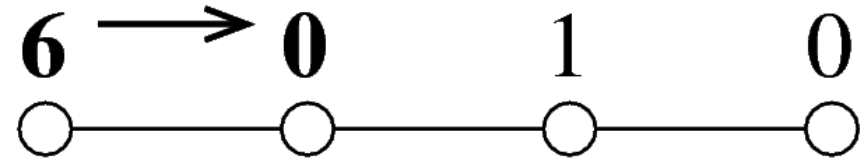
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Pebbling moves



Pebbling numbers

$\pi(G, D)$ is the number of pebbles required to ensure that D can be reached from any distribution of $\pi(G, D)$ pebbles.

If S is a **set** of distributions on G

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Optimal pebbling number:

$\pi^*(G, S)$ is the number of pebbles required in **some** distribution from which every $D \in S$ can be reached

Common pebbling numbers

$S_1(G)$: 1 pebble anywhere

$$\pi(G) = \pi(G, S_1(G))$$

(pebbling number)

$S_t(G)$: t pebbles on some vertex

$$\pi_t(G) = \pi(G, S_t(G))$$

(t -pebbling number)

δ_v : One pebble on v

$$\pi(G, v) = \pi(G, \delta_v)$$

$S(G, t)$ and $\pi(G, t)$

$S(G, t)$: all distributions with a total of t pebbles (anywhere on the graph)

Conjecture

$$\pi(G, S(G, t)) = \pi(G, S_t(G))$$

i.e. hardest-to-reach ***target*** configurations have all pebbles on one vertex

True for K_n , C_n , trees

Cover pebbling

$\Gamma(G)$: one pebble on every vertex

$$\gamma(G) = \pi(G, \Gamma(G))$$

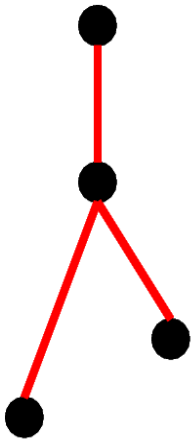
(*cover pebbling number*)

Sjöstrund: If $D(v) \geq 1$ for all vertices v , then there is a critical distribution with all pebbles on one vertex.

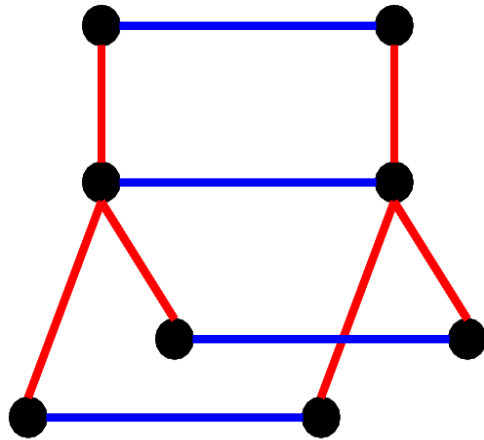
(A *critical* distribution has one pebble less than the required number and cannot reach some target distribution)

Cartesian products of graphs

K_2



G



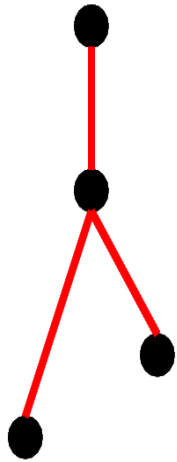
$G \times K_2$

Cartesian products of graphs

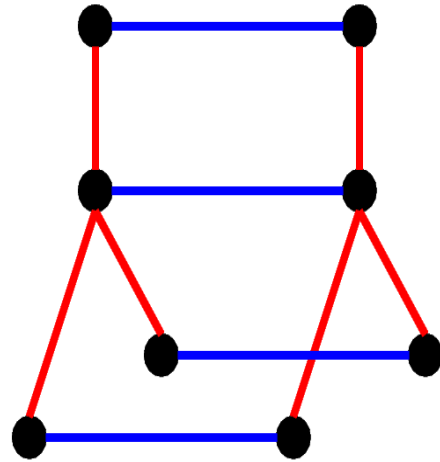
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K_3



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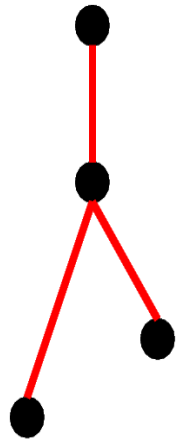
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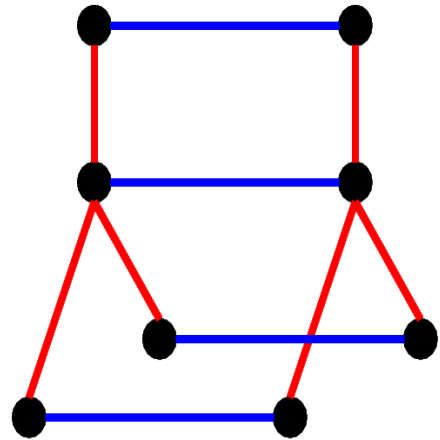
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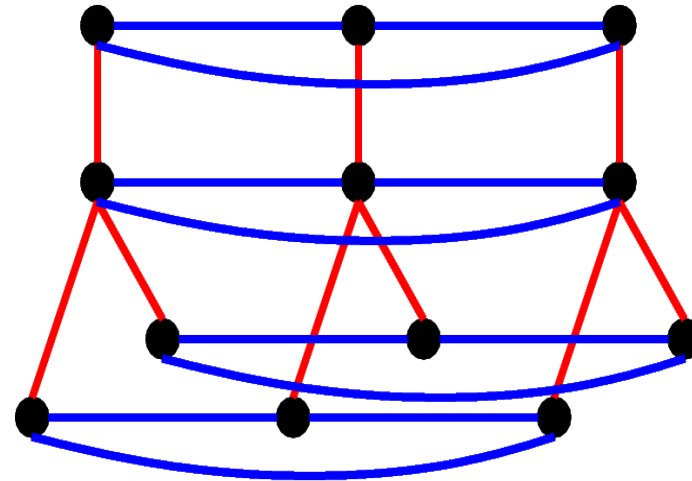
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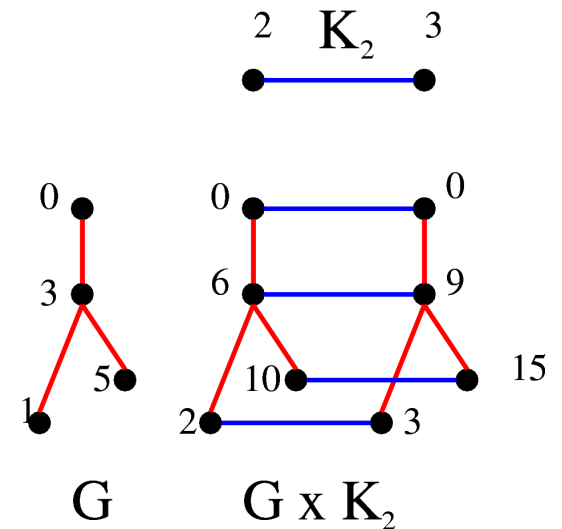
Products of distributions

Product of distributions:

D_1 on G ; D_2 on H

then $D_1 \cdot D_2$ on $G \times H$

$$D_1 \cdot D_2((v, w)) = D_1(v) D_2(w)$$



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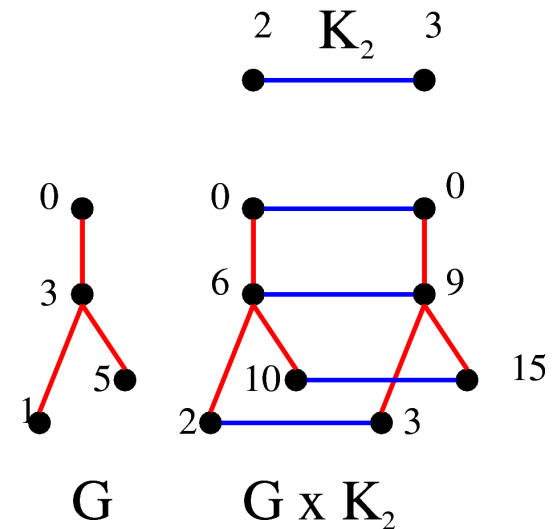
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Products of **sets** of distributions

S_1 a set of distros on G

S_2 a set of distros on H

$$S_1 \cdot S_2 = \{D_1 \cdot D_2 : D_1 \in S_1; D_2 \in S_2\}$$



Graham's Conjecture generalized

Graham's Conjecture:

$$\pi(G \times H) \leq \pi(G) \pi(H)$$

Generalization:

$$\pi(G \times H, S_1 \cdot S_2) \leq \pi(G, S_1) \pi(H, S_2)$$

Optimal pebbling

Observation (not obvious):

If we can get from D_1 to D_1' in G and from D_2 to D_2' in H , we can get from $D_1 \cdot D_2$ to $D_1 \cdot D_2'$ in $G \times H$ to $D_1' \cdot D_2'$ in $G \times H$

Conclusion (optimal pebbling):

$$\pi^*(G \times H, S_1 \cdot S_2) \leq \pi^*(G, S_1) \pi^*(H, S_2)$$

Graham's Conjecture holds for optimal pebbling in most general setting

Weighted graphs

Edges have weights w

Pebbling moves: remove w pebbles from one vertex, move 1 to adjacent vertex

$\pi(G)$, $\pi(G, D)$ and $\pi(G, S)$ still make sense

$G \times H$ also makes sense:

$$\text{wt}(\{(v, w), (v, w')\}) = \text{wt}(\{w, w'\})$$

$$\text{wt}(\{(v, w), (v', w)\}) = \text{wt}(\{v, v'\})$$

The Good news

Chung: Hypercubes ($K_2 \times K_2 \times \dots \times K_2$) satisfies Graham's Conjecture for any collection of weights on the edges

We can focus on complete graphs in most applications

The Bad News

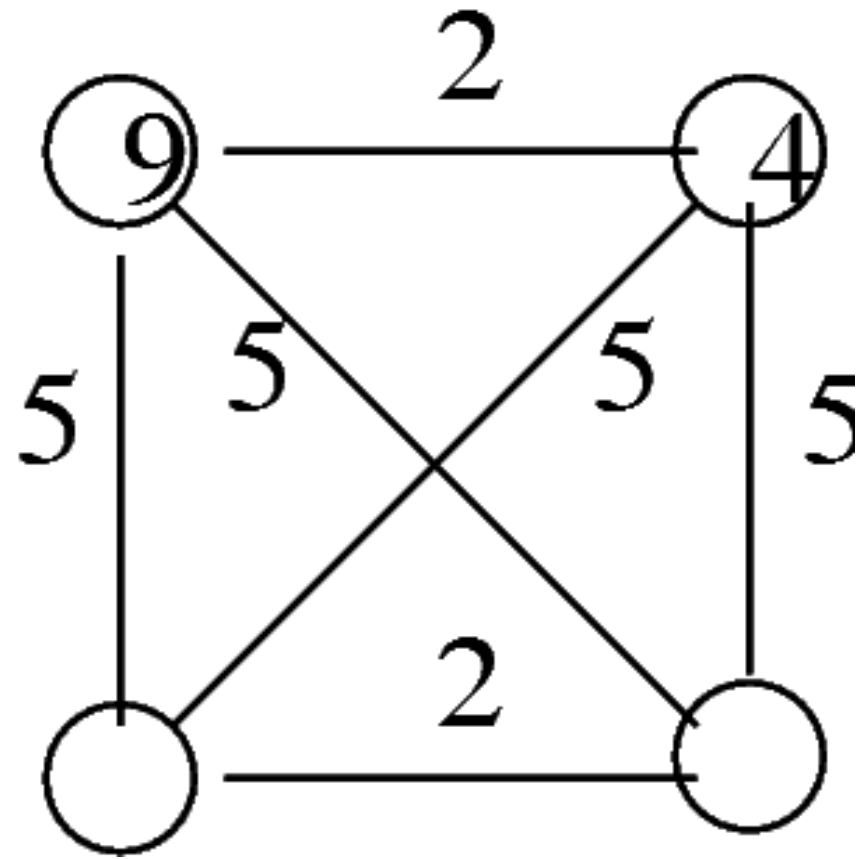
Complete graphs are
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The Bad News

Complete graphs are hard!

Sjöstrand's Theorem fails

13 pebbles on one vertex can cover K_4 , but...



Some specializations

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

Conjecture 2:

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

Clearly Conjecture 2 implies Conjecture 1

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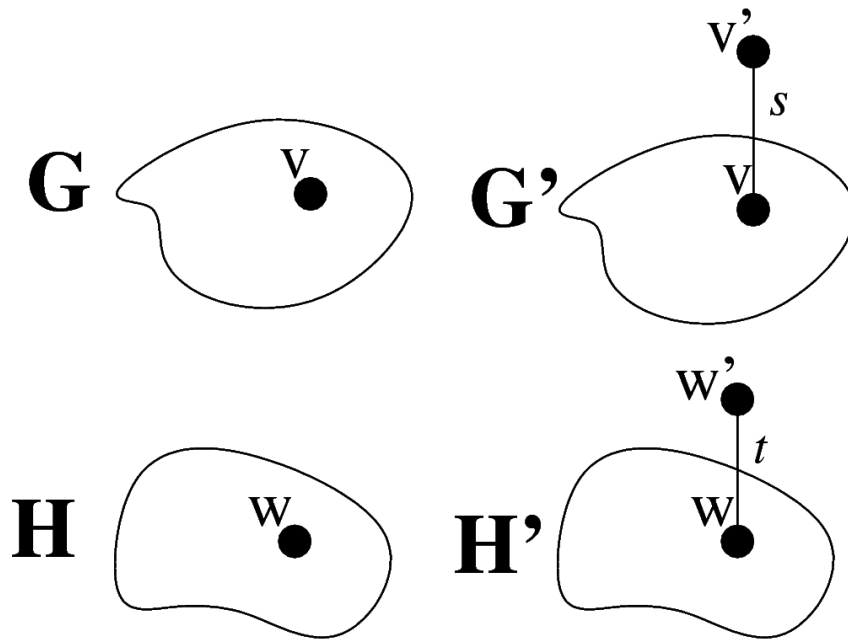
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These conjectures are equivalent on weighted graphs

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

implies

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$



$$\pi(G', v') = \pi_s(G, v)$$

$$\pi(H', w') = \pi_t(H, w)$$

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

implies

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

If st pebbles cannot be moved to (v, w) from D in $G \times H$, then (v', w') cannot be reached from D in $G' \times H'$ (delay moves onto $\{v'\} \times H'$ and $G' \times \{w'\}$ as long as possible)

$$\begin{aligned} \pi_{st}(G \times H, (v, w)) &\leq \pi(G' \times H', (v', w')) \\ &\leq \pi(G', v') \pi(H', w') = \pi_s(G, v) \pi_t(H, w) \end{aligned}$$

Implications for regular pebbling

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

equivalent to Conjecture 1':

$$\pi_{2^{ab}}(G \times H, (v, w)) \leq \pi_{2^a}(G, v) \pi_{2^b}(H, w)$$

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Conjecture 2:

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

equivalent to Conjecture 2':

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

when s and t are odd

Choosing a target

Observation: To reach an unoccupied vertex v in G , we need to put two pebbles on any neighbor of v .

We can choose the target neighbor

If S is a *set* of distributions on G , $\rho(G, S)$ is the number of pebbles needed to reach ***some*** distribution in S

Idea: Develop an induction argument to prove Graham's conjecture

Comparing pebbling numbers

$\forall D \in \mathcal{S} \forall D' \in \mathcal{S}(G, \pi(G, S))$

D is reachable from D' by a sequence of pebbling moves

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$\forall D \in S \forall D' \in S(G, \pi(G, S))$

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$\forall D \in S \exists D' \in S(G, \pi^*(G, S))$

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D is reachable from D' by a sequence of pebbling moves

$$\forall D \in S \exists D' \in S(G, \pi^*(G, S))$$

D is reachable from D' by a sequence of pebbling moves

$$\forall D' \in S(G, \rho(G, S)) \exists D \in S$$

D is reachable from D' by a sequence of pebbling moves

Properties of $\rho(G, S)$

$$\rho_1(G) = 1$$

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SURPRISE!

Graham's Conjecture fails!

Let H be the trivial graph: $S_H = \{2\delta_v\}$

$$\rho(G \times H, S_1(G) \cdot S_H) = \rho_2(G) = |V(G)| + 1$$

$$\rho(G, S_1(G)) = \rho_1(G) = 1$$

$$\rho(H, S_H) = 2$$